

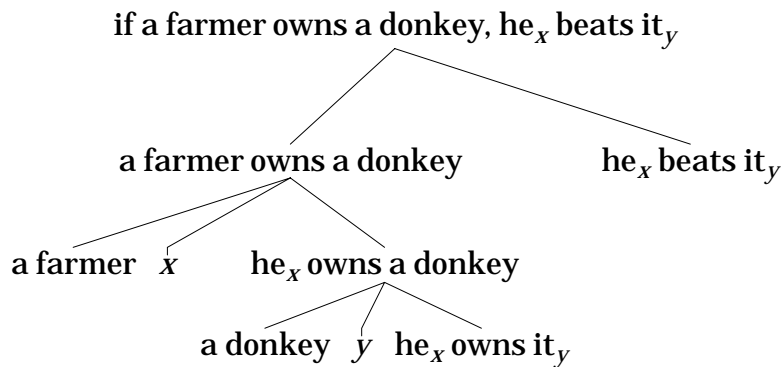
**1. From a classical point of view**

**a. Conditional donkeys**

(1) If a farmer owns a donkey, he beats it.

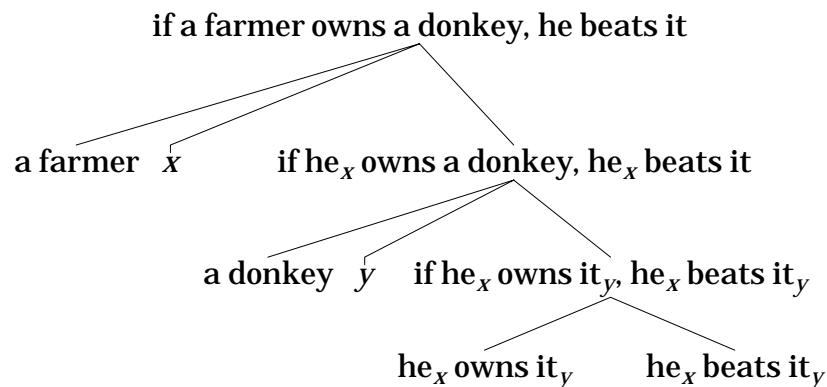
(2)  $(\forall x) (\forall y) [ [x \text{ is a farmer} \ \& \ y \text{ is a donkey} \ \& \ x \text{ owns } y] \rightarrow x \text{ beats } y ]$

(3)



(3')  $(\exists x) (\exists y) [ \underline{\text{farmer}'(x) \ \& \ \text{donkey}'(y) \ \& \ \text{own}'(x,y)} ] \Rightarrow \text{beat}'(x,y)$   
 ['If a farmer owns a donkey, then  $x$  owns  $y$ ']

(4)



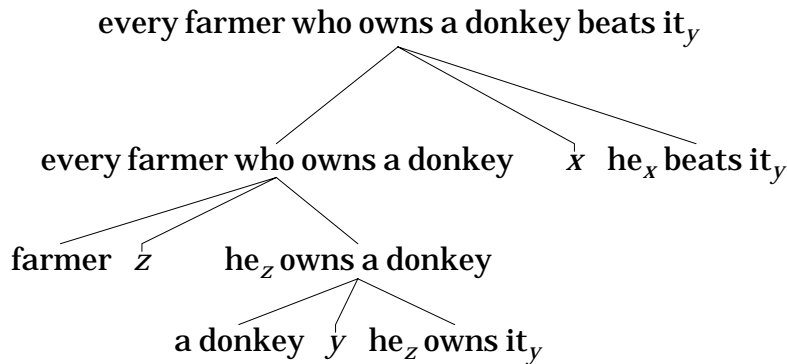
(4')  $(\exists x) (\exists y) [ \text{farmer}'(x) \ \& \ \text{donkey}'(y) \ \& \ (\text{own}'(x,y) \Rightarrow \text{beat}'(x,y)) ]$   
 ['A certain farmer beats a certain donkey, if he owns it.']

(5) ?If a farmer owns every/each donkey, he beats it.

**b. Relative donkeys**

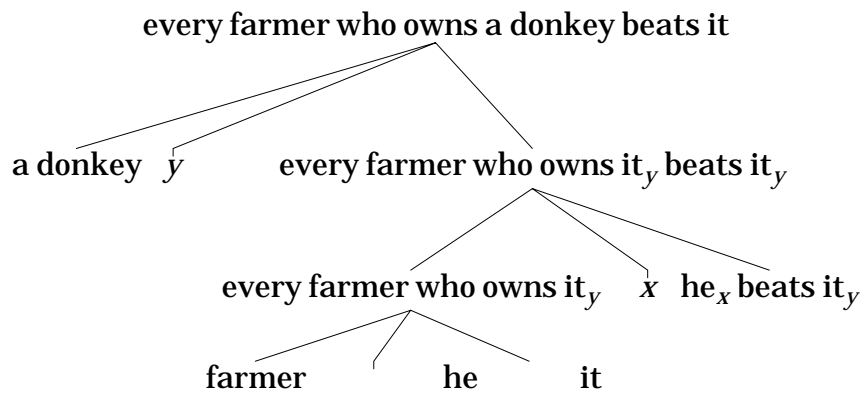
(6) Every farmer who owns a donkey beats it.

(7)



(7')  $(\forall x) ( [ \text{farmer}'(x) \ \& \ (\exists y) [ \text{donkey}'(y) \ \& \ \text{own}'(x,y) ] ] \rightarrow \text{beat}'(x,y) )$   
 ['Every farmer who owns a donkey beats y']

(8)



(8')  $(\exists y) [ \text{donkey}'(y) \ \& \ (\forall x) ( [ \text{farmer}'(x) \ \& \ \text{own}'(x,y) ] \rightarrow \text{beat}'(x,y) ) ]$   
 ['Every farmer who owns a certain donkey beats it']

(9) ?Every farmer who owns every/each donkey beats it.

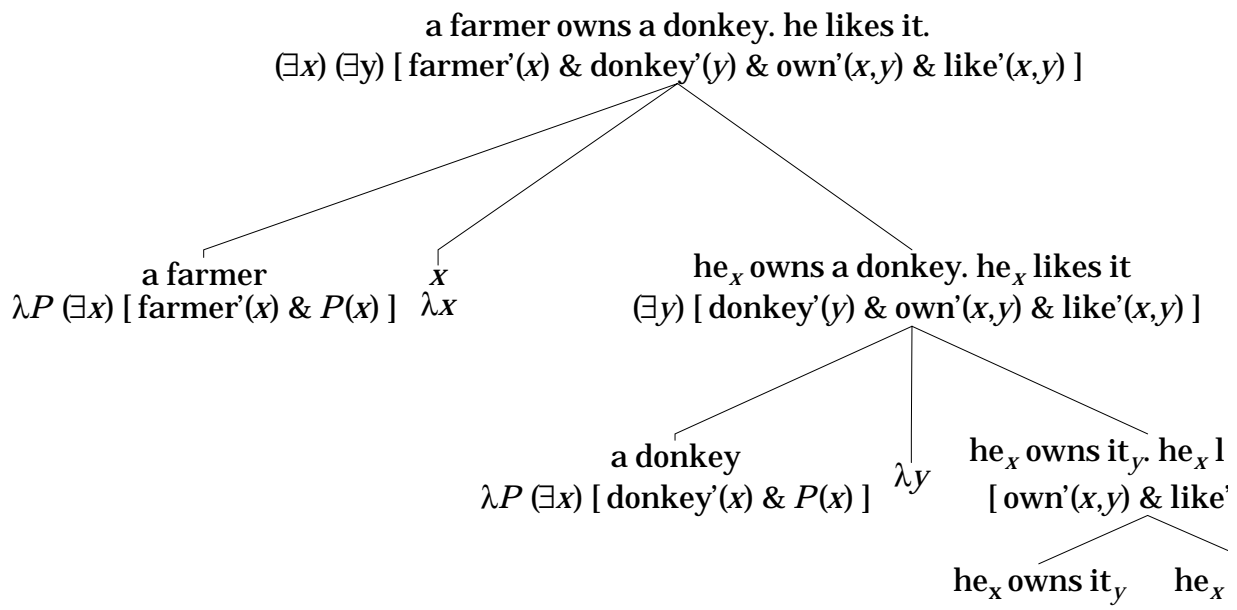
**2. Discourse Anaphora**

(10) A farmer owns a donkey. He likes it.

(11)

$$\underbrace{(\exists x) (\exists y)}_{?} \underbrace{[farmer'(x) \ \& \ donkey'(y) \ \& \ own'(x,y) \ \& \ like'(x,y)]}_{(a)} \underbrace{]}_{(b)}$$

(12)



(13) ?Every farmer owns a donkey. He likes it.

(14) The man who gave his paycheck to his wife was wiser than the one who gave it [i.e., **his** paycheck] to his mistress.

(15) (a) A boy owns a guinea-pig.

(b) He [i.e., **the** boy who owns guinea-pig] likes it [i.e., **the** guinea-pig that he, the boy who owns guinea-pig, owns].

(a')  $(\exists x) (\exists y) [ boy'(x) \ \& \ guinea-pig'(y) \ \& \ (own'(x,y)) ] .$

(b')  $like'((\lambda x) [ boy'(x) \ \& \ (\exists y) [ guinea-pig'(y) \ \& \ (own'(x,y)) ] ] ,$   
 $(\lambda y) [ guinea-pig'(y) \ \& \ (\exists x) [ boy'(x) \ \& \ (own'(x,y)) ] ] )$

- (16) (a) A farmer  $\left\{ \begin{array}{l} \text{rides on a bicycle} \\ \text{cycles} \end{array} \right\}$  .
- (b) It  $\left\{ \begin{array}{l} \text{[the bicycle that the farmer who rides on a bicycle rides on]} \\ \text{[the bicycle that the cycling farmer rides on]} \end{array} \right\}$  does  
not belong to him  $\left\{ \begin{array}{l} \text{[the farmer who rides on a bicycle]} \\ \text{[the cycling farmer]} \end{array} \right\}$  .
- (17) (a) A farmer rides on a bicycle.  
(b) It does not belong to him.  
(a')  $\lambda R (\exists x) (\exists y) [\text{farmer}'(x) \ \& \ \text{bicycle}'(y) \ \& \ \text{ride-on}'(x,y) \ \& \ \underline{R(x,y)}]$   
(b')  $\hat{x} \hat{y} [\neg \text{belong}'(y,x) \ \& \ R(x,y) ]$
- (18) (a) A farmer cycles.  
(b) It does not belong to him.  
(a')  $(\exists x) (\exists y) [\text{farmer}'(x) \ \& \ \text{bicycle}'(y) \ \& \ \text{ride-on}'(x,y) \ \& \ \underline{R(x)}]$   
(b')  $\hat{x} \hat{y} [\neg \text{belong}'(y,x) \ \& \ R(x) ]$
- (17') (a)  $\hat{x} \hat{y} [\text{farmer}'(x) \ \& \ \text{bicycle}'(y) \ \& \ \text{ride-on}'(x,y) ]$   
(b)  $\hat{x} \hat{y} \neg \text{belong}'(y,x)$   
(c)  $\hat{x} \hat{y} [\text{farmer}'(x) \ \& \ \text{bicycle}'(y) \ \& \ \text{ride-on}'(x,y) \ \& \ \neg \text{belong}'(y,x)]$
- (18') (a)  $\hat{x} (\exists y) [\text{farmer}'(x) \ \& \ \text{bicycle}'(y) \ \& \ \text{ride-on}'(x,y)]$   
(b)  $\hat{x} \hat{y} \neg \text{belong}'(y,x)$   
(c)  $\hat{x} \hat{y} [(\exists y) [\text{farmer}'(x) \ \& \ \text{bicycle}'(y) \ \& \ \text{ride-on}'(x,y) ] \ \& \ \neg \text{belong}'(y,x)]$
- (19) (a) A man loves a woman. He kisses her.  
(b) A man loves a woman. A man kisses her.

### 3. Adverbs of Quantification

- (20)
- If a farmer owns a donkey, he  $\left\{ \begin{array}{l} \text{always} \\ \text{sometimes} \\ \text{never} \\ \dots \end{array} \right\}$  beats it.

**(21)**

$$\left\{ \begin{array}{l} (\forall x) (\forall y) \\ (\exists x) (\exists y) \\ \neg (\exists x) (\exists y) \\ \dots \end{array} \right\} ([x \text{ is a farmer} \ \& \ y \text{ is a donkey} \ \& \ x \text{ owns } y] \quad \left\{ \begin{array}{l} \rightarrow \\ \& \\ \& \\ \dots \end{array} \right\} x \text{ beats } y)$$

**(21')**

$$\left\{ \begin{array}{l} \forall xy \\ \exists xy \\ \neg \exists xy \\ \dots \end{array} \right\} ([x \text{ is a farmer} \ \& \ y \text{ is a donkey} \ \& \ x \text{ owns } y] \quad \left\{ \begin{array}{l} \rightarrow \\ \& \\ \& \\ \dots \end{array} \right\} x \text{ beats } y)$$

**(21'')**

$$\left\{ \begin{array}{l} \forall \\ \exists \\ \neg \exists \\ \dots \end{array} \right\} (\hat{x} \hat{y} ([x \text{ is a farmer} \ \& \ y \text{ is a donkey} \ \& \ x \text{ owns } y] \quad \left\{ \begin{array}{l} \rightarrow \\ \& \\ \& \\ \dots \end{array} \right\} x \text{ beats } y))$$

**(22)** If a farmer owns a donkey, he usually beats it.

**(23)**

$$\left\{ \begin{array}{l} \forall \\ \exists \\ \neg \exists \\ \text{MOST} \\ \dots \end{array} \right\} (\hat{x} \hat{y} \ x \text{ is a farmer} \ \& \ y \text{ is a donkey} \ \& \ x \text{ owns } y, \quad \hat{x} \hat{y} \ x \text{ beats } y)$$

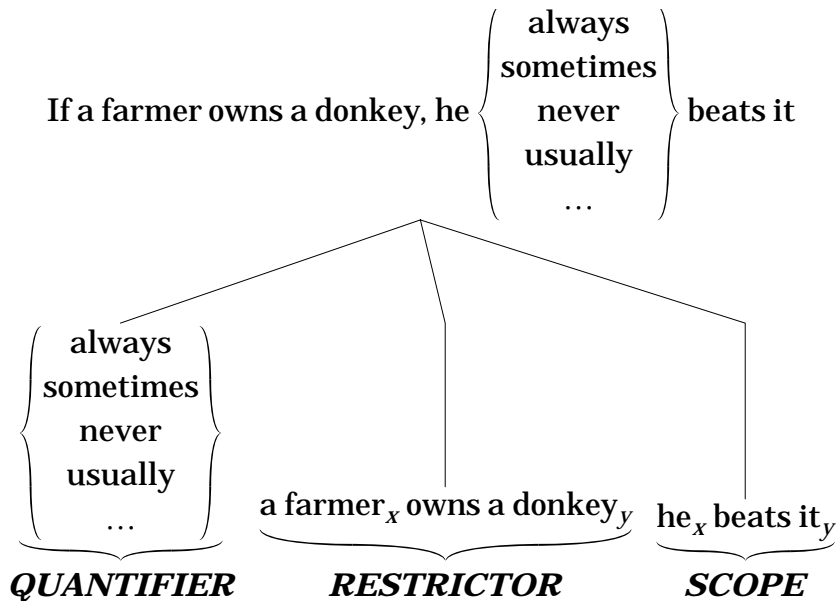
**(24)**

$$\text{If a boy draws a picture of a girl, he} \quad \left\{ \begin{array}{l} \text{always} \\ \text{sometimes} \\ \text{never} \\ \text{usually} \\ \dots \end{array} \right\} \text{ gives it to her.}$$

**(25)** If  $\varphi(a N_1, \dots, a N_n)$ , [then] ADV  $\psi(it_1, \dots, it_n) \mapsto$

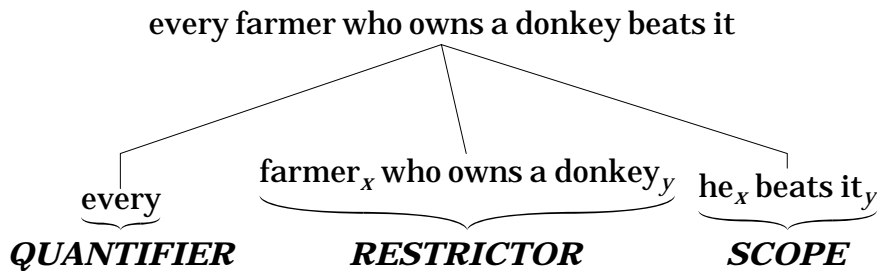
$$\text{ADV}'(\hat{x}_1, \dots, \hat{x}_n [x_1 \text{ is a } N_1 \ \& \ \dots \ \& \ x_n \text{ is a } N_n \ \& \ \varphi'(x_1, \dots, x_n)], \hat{x}_1, \dots, \hat{x}_n \psi'(x_1, \dots, x_n))$$

(26)



- (27)  $\text{a farmer}_x \text{ owns a donkey}_y \mapsto [ \text{farmer}'(x) \ \& \ \text{donkey}'(y) \ \& \ \text{own}'(x,y) ]$   
 $\text{a farmer}_x \mapsto \text{farmer}'(x)$  ('Indefinites as variables')  
 $\text{a farmer owns a donkey} \mapsto [ (\text{farmer}' \times \text{donkey}') \cap \text{own}' ]$   
 $\text{a farmer}_x \mapsto \text{farmer}'$  ('Indefinites as properties')

(28)



#### 4. Asymmetries

- (29) Most farmers who own a donkey beat it.  
 (29')  $\text{MOST}(\hat{x} \hat{y} \text{ farmer}'(x) \ \& \ \text{donkey}'(y) \ \& \ \text{own}'(x,y), \quad \hat{x} \hat{y} \text{ beat}'(x,y))$

- (30) Every person who has a dime will put it in the meter.
- (30')  $\forall (\hat{x} \hat{y} \text{ person}'(x) \ \& \ \text{dime}'(x) \ \& \ \text{have}'(x,y) \ \& \ \text{put-in-the-meter}'(x,y))$   
 $[= (\forall x) (\forall y) ( [\text{person}'(x) \ \& \ \text{dime}'(x) \ \& \ \text{have}'(x,y)] \rightarrow$   
 $\text{put-in-the-meter}'(x,y))$   
 $=$   
 $(\forall x) ( [\text{person}'(x) \rightarrow$   
 $(\forall y) ( [ \text{dime}'(x) \ \& \ \text{have}'(x,y)] \rightarrow \text{put-in-the-meter}'(x,y) ) ] ) ]$
- (30'')  $(\forall x) ( (\exists y) [\text{person}'(x) \ \& \ \text{dime}'(x) \ \& \ \text{have}'(x,y)] \rightarrow$   
 $(\exists y) [\text{person}'(x) \ \& \ \text{dime}'(x) \ \& \ \text{have}'(x,y) \ \& \ \text{put-in-the-meter}'(x,y) ] )$
- (31)  $\text{every}_{\text{universal}}^n =$   
 $\lambda R \lambda S (\forall x) (\forall y_2) \dots (\forall y_n) [ R(x,y_2,\dots,y_n) \rightarrow S(x,y_2,\dots,y_n) ]$
- (32)  $\text{every}_{\text{existential}}^n =$   
 $\lambda R \lambda S (\forall x) [ (\exists y_2) \dots (\exists y_n) R(x,y_2,\dots,y_n) \rightarrow$   
 $(\exists y_2) \dots (\exists y_n) [ R(x,y_2,\dots,y_n) \ \& \ S(x,y_2,\dots,y_n) ] ]$
- (33) Most persons who have a dime will put it in the meter.
- (33')  $\text{MOST} ( \hat{x} (\exists y) [\text{person}'(x) \ \& \ \text{dime}'(x) \ \& \ \text{have}'(x,y)] ,$   
 $\hat{x} (\exists y) [\text{person}'(x) \ \& \ \text{dime}'(x) \ \& \ \text{have}'(x,y) \ \& \ \text{will-put-in-the-meter}'(x,y)] ) )$
- (34)  $\text{EXISTENTIAL}(Q^n) =$   
 $\lambda R \lambda S (Qx) ((\exists y_2) \dots (\exists y_n) R(x,y_2,\dots,y_n) ,$   
 $(\exists y_2) \dots (\exists y_n) [ R(x,y_2,\dots,y_n) \ \& \ S(x,y_1,\dots,y_n) ] )$   
 $[= \lambda R \lambda S Q(\hat{x} (\exists y_2) \dots (\exists y_n) R(x,y_2,\dots,y_n) ,$   
 $\hat{x} (\exists y_2) \dots (\exists y_n) [ R(x,y_2,\dots,y_n) \ \& \ S(x,y_1,\dots,y_n) ] ) ]$
- (35) Every farmer who owns a donkey beats it.
- (35')  $(\forall x) ( (\exists y) [\text{farmer}'(x) \ \& \ \text{donkey}'(x) \ \& \ \text{have}'(x,y)] \rightarrow$   
 $(\exists y) [\text{farmer}'(x) \ \& \ \text{donkey}'(x) \ \& \ \text{have}'(x,y) \ \& \ \text{beat}'(x,y) ] )$
- (36) Most people that owned a slave also owned his offspring.
- (36')  $\text{MOST}( \hat{x} \hat{y} [\text{person}'(x) \ \& \ \text{slave}'(y) \ \& \ \text{own}'(x,y)],$   
 $\hat{x} \hat{y} [\text{person}'(x) \ \& \ \text{slave}'(y) \ \& \ \text{own}'(x,y's \ \text{offspring}) ] )$
- (37)  $\text{MOST}( \hat{x} (\exists y) [\text{person}'(x) \ \& \ \text{slave}'(y) \ \& \ \text{own}'(x,y)],$   
 $\hat{x} (\forall y) [ [\text{person}'(x) \ \& \ \text{slave}'(y) \ \& \ \text{own}'(x,y)] \rightarrow \text{own}'(x,y's \ \text{offspring}) ] )$

- (38) UNIVERSAL( $Q^n$ ) =  
 $\lambda R \lambda S (Qx) ((\exists y_2) \dots (\exists y_n) R(x, y_2, \dots, y_n) ,$   
 $(\forall y_2) \dots (\forall y_n) [R(x, y_2, \dots, y_n) \rightarrow S(x, y_1, \dots, y_n) ]$
- (39) If a farmer owns a donkey he is usually rich.  
(39') (MOST  $\hat{x} \hat{y}$ ) ( [farmer'(x) & donkey'(y) & own'(x,y)] , rich'(x) ] )  
(39'') (MOST  $\hat{x}$ ) (  $(\exists y)$  [farmer'(x) & donkey'(y) & own'(x,y), rich'(x)]  
(= (MOST  $\hat{x}$ ) (  $(\exists y)$  [farmer'(x) & donkey'(y) & own'(x,y),  
 $(\exists y)$  [farmer'(x) & donkey'(y) & own'(x,y) & rich'(x)] ) )  
=  
(EXISTENTIAL(MOST)  $\hat{x}$ ) ([farmer'(x) & donkey'(y) & own'(x,y)], rich'(x) ) )
- (40) If a DRUMMER lives in an apartment complex, it is usually half empty.  
(40') (MOST  $\hat{y}$ ) (  $(\exists x)$  [apartment-complex'(x) & drummer'(y) & live-in'(y,x)],  
 $(\exists y)$  [apartment-complex'(x) & drummer'(y) & live-in'(y,x) & half-empty'(y) ] )  
['The majority of apartment complexes with a drummer in them are half empty.']
- (41) If a drummer lives in an APARTMENT COMPLEX, it is usually half empty.  
(41') (MOST  $\hat{x}$ ) (  $(\exists y)$  [apartment-complex'(x) & drummer'(y) & live-in'(y,x)],  
 $(\exists y)$  [apartment-complex'(x) & drummer'(y) & live-in'(y,x) & half-empty'(y) ] )  
['The majority of drummers living in apartment complexes live in half empty apartment complexes']
- 5. Uniqueness**
- (42) Every farmer who owns a donkey beats it.  
(42')  $(\forall x) ( [ \text{farmer}'(x) \ \& \ (\exists y) [ \text{donkey}'(y) \ \& \ \text{own}'(x,y) ] ] ,$   
 $\text{beat}'(x, \underline{1}y) [ \text{donkey}'(y) \ \& \ \text{own}'(x,y) ] )$   
['Every farmer beats *the* donkey he owns']
- (43) Every woman who bought a sage plant bought eight others along with it.  
(44) No parent with a teenage son lends him the car.
- (45) If a woman buys a sage plant here, she always buys eight others along with it.  
(46) If a woman has a teenage son, she never lends him the car.



### 1. Indirect interpretation

(1)

N	→	L	⇒	M	⊇	Γ
natural language	<i>translation algorithm</i>	formal language (logic)	<i>inter- pretation</i>	class of models	<i>restricted by meaning postulates</i>	‘good’ models

$\Gamma = \{m \in M \mid M \models \Pi\} = \{m \in M \mid \forall \phi \in \Pi: M \models \phi\}$   
 $m = (F, i, c, g)$

### 2. Kinds of postulates

(2)  $\phi = \phi(c_1, \dots, c_n)$

(3)  $be' = \hat{P} \hat{x} P\{\hat{y} \ x = y\}$

(4)  $every' = \hat{P} \hat{Q} (\forall x) [P\{x\} \rightarrow Q\{x\}]$

(5a) **John is slowly eating a banana.**  
 $\therefore$  **John is eating a banana.**

(5b)  $(\forall P) (\forall x) [slowly'(x) \rightarrow P\{x\}]$

(6a) **John finds a banana.**  
 $\therefore$  **There is a banana.**

(6b) **Mary found a groundhog.**  
**All woodchucks are groundhogs.**  
**All groundhogs are woodchucks.**  
 $\therefore$  **John found a woodchuck.**

(6c)  $(\exists R) find = \hat{P} \hat{x} P\{\hat{y} \ R\{x, y\}\}$

(7a) **Mary loves John.**                      **Everyone loves John.**  
 $\therefore$  **Someone loves John.**                       $\therefore$  **Mary loves John.**

(7b)  $(\exists x) Mary' = \hat{P} P\{x\}$

(8) **kill = cause'(die)**

(9)  $seek' = \hat{F} \hat{x} try'(x, P\{\hat{y} \ find(x, y)\})$

(10a)  $Mary' = \hat{P} P(m)$

(10b)  $Mary' = \lambda i \lambda P P(m(i))$

(10c)  $Mary' = \lambda i \lambda P P(m)$

(10)  $(\forall x) [ \textit{bachelor}'(x) \rightarrow \neg \textit{married}'(x) ]$

(11)  $no' = \hat{P} \hat{Q} \textit{every}'(P) (\hat{x} \neg Q\{x\})$

(12)  $(\forall x) [ \textit{pilot}'(x) \rightarrow (\exists y) [ \textit{plane}'(y) \ \& \ \textit{fly}'(x,y) ] ]$

(13)  $(\forall y) [ \textit{plane}'(y) \rightarrow (\exists x) [ \textit{pilot}'(x) \ \& \ \textit{fly}'(x,y) ] ]$

### 3. Problems with postulates

(14a)        *John saw nobody smile.*  
 $\therefore$         *There was nobody that John saw smile.*

(14b)        *There was nobody that John saw smile.*  
 $\therefore$         *John saw nobody smile.*

(14c)  $(\forall x) (\forall Q) (\forall P) [ \textit{see}'(x, \hat{Q}(\hat{y} P\{y\})) \leftrightarrow Q(\hat{y} \textit{see}'(x, \hat{P}\{y\})) ]$   
 $[ = \dots [ \textit{see}'(x, \hat{(Qy) P\{y\}}) \leftrightarrow (Qy) \textit{see}'(x, \hat{P}\{y\}) ] ]$

(15) *John opened the drawer. Mary closed it again<sub>2</sub>.*

(16) *It rained again<sub>1</sub>.*

(17) *John opened the drawer. Mary had closed it again<sub>2</sub>.*

(18)  $(\forall x) (\forall P) (\forall p)$   
 $[ \textit{again}_2(\wedge \textit{CAUSE}(\wedge P\{x\}), \wedge \textit{BECOME}(p))$   
 $\leftrightarrow \textit{CAUSE}(\wedge P\{x\}, \wedge \textit{again}_1(\wedge \textit{BECOME}(p))) ]$

(19)  $\square (\forall x) (\forall P) (\forall p)$   
 $[ \textit{CAUSE}(\wedge \textit{press-the-button}(x), \wedge \textit{BECOME}(\textit{water}))$   
 $\leftrightarrow \textit{CAUSE}(\wedge \textit{press-the-button}(x), \wedge \textit{BECOME}(\textit{cold-water})) ]$

(19')  $[ \textit{again}_2(\wedge \textit{CAUSE}(\wedge \textit{press-the-button}(p)), \wedge \textit{BECOME}(\textit{water})) \leftrightarrow$   
 $\textit{CAUSE}(\wedge \textit{press-the-button}(p), \wedge \textit{again}_1(\wedge \textit{BECOME}(\textit{cold-water}))) ]$