## Compositionality

## and

## Ontological Commitment

Thomas Ede Zimmermann, Goethe University, Frankfurt SPE 5, University of Turin, July 25-27, 2012
0. Intro
I. The compositional enterprise
2. External extensions
3. Internal Extensions
4. Intensions and Fregean compositionality
5. Outro

## 0. Intro

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To be is to be the value of a (bound) variable.

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Quine (1961)

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$(\forall x)[M(x) \rightarrow(\exists y)[W(y) \& L(x, y)]]$

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ranging over...

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$(\lambda Y . \lambda X . Y \subseteq X)(M)(\lambda x .(\lambda Y . \lambda X . Y \notin X)(W)(\{y \mid x L y\}))$

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To be is to be the value of a (bound) variable. Quine (1961)
Ontological overkill:
Every man loves a woman.
Non-compositional analysis (formalisation)
$(\nexists x)[M(x) \rightarrow(\exists y)[W(y) \& L(x, y)]]$ $\xrightarrow{ }$ individuals ranging over...

Compositional analysis (indirect interpretation)
$(\lambda M) X . Y \subseteq X)(M)(\lambda x .(X Y) X X . Y \nsubseteq X)(W)(\{y \mid x L y\}))$

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To be is to be the value of a (bound) variable. Quine (1961)
Ontological overkill:
Every man loves a woman.
Non-compositional analysis (formalisation)
$(\forall x)[M(x) \rightarrow(\exists y)[W(y) \& L(x, y)]]$ individuals ranging over...

Compositional sets of individuals
$(\lambda Y) \lambda X . Y \subseteq X)(M)(\lambda \times \cdot(\lambda Y), X X) \times X)(W)(\{y \mid x L y\}))$

## 0. Intro

To be is to be the value of a (bound) variable.
Ontological overkill:
Every man loves a woman.
Non-compositional analysis (formalisation)
$(\notin x)[M(x) \rightarrow(\exists y)[W(y) \& L(x, y)]]$ $\xrightarrow{ }$ individuals ranging over...

Compositional sets of individuals
$(\lambda M) X X . Y \subseteq X)(M)(\lambda x .(X Y) X X . Y \nVdash X)(W)(\{y \mid x L y\}))$

## 0. Intro

To be is to be the referent of a term.
Ontological overkill:
Every man loves a woman.
Non-compositional analysis (formalisation)
$(\forall x)[M(x) \rightarrow(\exists y)[W(y) \& L(x, y)]]$ ranging over...
sets of individuals
Compositional (indirect interpretation)
$(\lambda Y . \lambda X . Y \subseteq X)(M)(\lambda x .(\lambda Y . \lambda X . Y \notin X)(W)(\{y \mid x L y\}))$

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To be is to be the referent of a term.
Ontological overkill:
Every man seeks a woman.
Non-compositional analysis (formalisation)
$(\forall x)[M(x) \rightarrow(\exists y)[W(y) \& L(x, y)]]$ ranging over...
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Compositional (indirect interpretation)
$(\lambda Y . \lambda X . Y \subseteq X)(M)(\lambda x .(\lambda Y . \lambda X . Y \notin X)(W)(\{y \mid x L y\}))$

## 0. Intro

To be is to be the referent of a term.
Ontological overkill:
Every man seeks a woman.
Non-compositional analysis (formalisation)

$$
(\forall x)[M(x) \rightarrow \square_{x} \underbrace{(\exists y)}_{\text {set of worlds (proposition) }}[W(y) \& L(x, y)]]
$$

denoting ...
$\frac{\text { Compositional sets of individuals }}{(\lambda Y . \lambda X . Y \subseteq X)(M)(\lambda \times .(\lambda Y . \lambda X . Y \propto X)(W)(\{y \mid x L y\}))}$

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To be is to be the referent of a term.
Ontological overkill:
Every man seeks a woman.
Non-compositional analysis (formalisation)

$$
(\forall x)[M(x) \rightarrow \Phi_{x} \underbrace{(\exists y)}_{\text {set of worlds (proposition) }}[W(y) \& L(x, y)]]
$$

denoting ...
function from worlds to sets of sets of individuals
Compositional analysis (indirect interpretation)
$(\lambda Y . \lambda X . Y \subseteq X)(M)\left(\lambda x . S\left(x,^{\wedge}(\lambda Y . \lambda X . Y \nless X)(W)\right)\right)$

## 0. Intro

I. The compositional enterprise
2. External extensions
3. Internal extensions
4. Intensions and Fregean compositionality
5. Outro
0. Intro
I. The compositional enterprise
2. External extensions
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5. Outro

## I. The compositional enterprise

## I. The compositional enterprise

## I. The compositional enterprise

## START

## I. The compositional enterprise

## START

EXPRESSIONS

## I. The compositional enterprise

## START

EXPRESSIONS

## I. The compositional enterprise

## START

EXPRESSIONS have

$$
\text { semantic values }\langle X\rangle, 《 X\rangle, \ldots
$$

## I. The compositional enterprise

## START

EXPRESSIONS have
semantic values $\langle X\rangle, 《 X\rangle, \ldots$
corresponding to their (communicative,...) functions: reference (potential), informational content,...

## I. The compositional enterprise

## START

Some
EXPRESSIONS have external semantic values $\langle X\rangle, 《 X\rangle, \ldots$ corresponding to their (communicative,...) functions: reference (potential), informational content,...

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Some
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GOAL

## I．The compositional enterprise

## START

Some

EXPRESSIONS have external semantic values $\langle X\rangle, 《 X\rangle, \ldots$ corresponding to their（communicative，．．．）functions： reference（potential），informational content，．．．

## GOAL

All
EXPRESSIONS have external（or internal）semantic values 〈X〉，《X》，．．． corresponding to（the contributions they make to） the functions of expressions in which they occur

## I．The compositional enterprise

## START

Some

EXPRESSIONS have external semantic values $\langle X\rangle, 《 X\rangle, \ldots$ corresponding to their（communicative，．．．）functions： reference（potential），informational content，．．．

## GOAL

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EXPRESSIONS have external（or internal）semantic values 〈X〉，《X》，．．． corresponding to（the contributions they make to） the functions of expressions in which they occur

## I. The compositional enterprise GOAL <br> All

EXPRESSIONS have
external (or internal) semantic values $\langle\mathrm{X}\rangle, 《 \mathrm{X}\rangle, \ldots$ corresponding to (the contributions they make to) the functions of expressions in which they occur

## I．The compositional enterprise GOAL <br> All

EXPRESSIONS have
external（or internal）semantic values 〈X〉，《X》，．．． corresponding to（the contributions ${ }^{*}$ they make to） the functions of expressions in which they occur
＊that are compositional：

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EXPRESSIONS have external（or internal）semantic values 〈X〉，《X》，．．． corresponding to（the contributions＊they make to） the functions of expressions in which they occur
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＊that are compositional：

$$
\left\langle\begin{array}{ll}
X+Y \\
X & Y \\
\Delta & \Delta
\end{array}\right\rangle=
$$

## I．The compositional enterprise GOAL <br> All

EXPRESSIONS have external（or internal）semantic values 〈X〉，《X》，．．． corresponding to（the contributions＊they make to） the functions of expressions in which they occur
＊that are compositional：

$$
\left\langle\begin{array}{cc}
X+Y \\
X & Y \\
\Delta & \Delta
\end{array}\right\rangle=\left\langle\begin{array}{c}
X \\
\Delta
\end{array}\right\rangle \oplus\left\langle\begin{array}{l}
Y \\
\Delta
\end{array}\right\rangle
$$

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I. The compositional enterprise
6. External extensions
7. Internal extensions
8. Intensions and Fregean compositionality
9. Outro

## 2. External extensions

## 2. External extensions

## 2. External extensions

|st approach

## 2. External extensions

Ist approach
Frege (1892)

## 2. External extensions

Ist approach
Frege (1892)
START

## 2. External extensions

Ist approach
START

EXPRESSIONS have external
semantic values $\langle X\rangle, 《 X\rangle, \ldots$ corresponding to their (communicative,...) functions: reference (potential), informational content,...

## 2. External extensions

Ist approach
START

EXPRESSIONS have
external
semantic values $\langle X\rangle, 《 X\rangle, \ldots$
corresponding to their (communicative,...) functions: reference (potential), informational content,...

## 2. External extensions

Ist approach
START with
REFERENTIAL
EXPRESSIONS have
external
semantic values $\langle X\rangle, 《 X\rangle, \ldots$
corresponding to their (communicative,...) functions: reference (potential), informational content,...

## 2. External extensions

Ist approach
START
REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function

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Ist approach
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REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function

$\left\langle\right.$| NP |  |
| :---: | :---: |
| Det | N |
| the | King |$\rangle$

## 2. External extensions

Ist approach
START
REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function

$\left\langle\right.$| $N P$ |  |
| :---: | :---: |
| Det | $N$ |
| the | King |$\rangle=$

## 2. External extensions

Ist approach
START
REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function

$\left\langle\right.$| $N P$ |  |
| :---: | :---: |
| Det | $N$ |
| the | King |$\rangle=$ Elvis

## 2. External extensions

Ist approach
START
REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function $\left\langle\begin{array}{c}\mathrm{NP} \\ \text { Elvis }\end{array}\right\rangle=$ Elvis

## 2. External extensions

Ist approach
START
REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function
$\left\langle\begin{array}{c}\mathrm{NP} \\ \text { Elvis }\end{array}\right\rangle=$ Elvis

+ (somewhat mysteriously)


## 2. External extensions

Ist approach
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REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function


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Ist approach
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REFERENTIAL
EXPRESSIONSget their referents as external semantic values $\langle\mathrm{X}\rangle$ corresponding to their referential function

$=1$

## 2. External extensions

2nd approach
START
TERMS (= Names + Descriptions): as before PREDICATES get their satisfiers as external semantic values $\langle P\rangle$ corresponding to their multiple reference

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2nd approach
START
TERMS (= Names + Descriptions): as before PREDICATES get their satisfiers as external semantic values $\langle P\rangle$ corresponding to their multiple reference

〈is dead〉

## 2. External extensions

2nd approach
START
TERMS (= Names + Descriptions): as before PREDICATES get their satisfiers as external semantic values $\langle P\rangle$ corresponding to their multiple reference

〈is dead> $\quad=\{(\mathrm{x}) \mid \mathrm{x}$ is dead $\}$

## 2．External extensions

2nd approach
START
TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfiers as external semantic values 〈P〉 corresponding to their multiple reference

〈is dead〉

$$
=\{(x) \mid x \text { is dead }\}
$$

〈is watching〉

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfiers as external semantic values 〈 P$\rangle$ corresponding to their multiple reference

〈is dead〉 $\quad=\{(x) \mid x$ is dead $\}$
〈is watching〉 $\quad=\{(\mathrm{x}, \mathrm{y}) \mid \mathrm{x}$ is watching y$\}$

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfiers as external semantic values 〈P〉 corresponding to their multiple reference

〈is dead〉 $\quad=\{(\mathrm{x}) \mid \mathrm{x}$ is dead $\}$
〈is watching〉 $=\{(\mathrm{x}, \mathrm{y}) \mid \mathrm{x}$ is watching y$\}$
〈is showing〉

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfiers as external semantic values $\langle P\rangle$ corresponding to their multiple reference

〈is dead〉 $\quad=\{(\mathrm{x}) \mid \mathrm{x}$ is dead $\}$
〈is watching〉 $=\{(x, y) \mid x$ is watching $y\}$
〈is showing〉
$=\{(x, y, z) \mid x$ is showing $y$ to $z\}$

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfiers as external semantic values $\langle P\rangle$ corresponding to their multiple reference valency

```
| 〈is dead\rangle = {(x)| }\textrm{x}\mathrm{ is dead}
2 〈is watching\rangle = {(x,y)| x is watching y}
3〈is showing\rangle ={(x,y,z)|x}\mathrm{ is showing }\textrm{y}\mathrm{ to z}
```


## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfiers as external semantic values $\langle P\rangle$ corresponding to their multiple reference valency ＝

I 〈is dead〉 $\quad=\{(x) \mid x$ is dead $\}$
2 〈is watching〉 $=\{(x, y) \mid x$ is watching $y\}$
3 〈is showing〉 $\quad=\{(x, y, z) \mid x$ is showing $y$ to $z\}$

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfiers as external semantic values $\langle P\rangle$ corresponding to their multiple reference
$=\{(x, y) \mid x$ is watching $y\}$
3 〈is showing〉
$=\{(x, y, z) \mid x$ is showing $y$ to $z\}$

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfies as external semantic values $\langle P\rangle$ corresponding to their multiple reference

〈Elvis is dead
I 〈 is dead 〉 $\quad=\{(x) \mid x$ is dead $\}$
2 〈is watching
$=\{(x, y) \mid x$ is watching $y\}$
3 〈is showing
$=\{(x, y, z) \mid x$ is showing $y$ to $z\}$

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfies as external semantic values $\langle P\rangle$ corresponding to their multiple reference

0 〈Elvis is dead
I 〈 is dead 〉 $\quad=\{(x) \mid x$ is dead $\}$
2 〈is watching
$=\{(x, y) \mid x$ is watching $y\}$
3 〈is showing
$=\{(x, y, z) \mid x$ is showing $y$ to $z\}$

## 2．External extensions

2nd approach

TERMS（＝Names＋Descriptions）：as before PREDICATES get their satisfies as external semantic values $\langle P\rangle$ corresponding to their multiple reference
valency
$=$
0 〈Elvis is dead
I 〈 is dead 〉 $\quad=\{(x) \mid x$ is dead $\}$
2 〈is watching
$=\{(x, y) \mid x$ is watching $y\}$
$=\{(x, y, z) \mid x$ is showing $y$ to $z\}$
－arity

3 〈is showing

## 2. External extensions

2nd approach

TERMS (= Names + Descriptions): as before PREDICATES get their satisfies as external semantic values $\langle P\rangle$ corresponding to their multiple reference
$0\langle$ Elvis is dead $\quad=\{() \mid$ Elvis is dead $\}$
$=\{(x) \mid x$ is dead $\}$
$=\{(x, y) \mid x$ is watching $y\}$
3 〈is showing
$=\{(x, y, z) \mid x$ is showing $y$ to $z\}$

## 2. External extensions

Comparison

# 2. External extensions 

Comparison

Names
Descriptions
Nouns

Verbs

Sentences
Determiners

## 2. External extensions

Comparison
Frege
Carnap

Names
Descriptions
Nouns
Verbs
Sentences
Determiners

## 2. External extensions

Comparison

Frege

Names
Descriptions individuals
Nouns
Verbs
Sentences
Determiners
individuals
—
-
truth values
-
$\qquad$

## Carnap

individuals individuals sets relations
truth values -
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I. The compositional enterprise
2. External extensions

## 3. Internal extensions

4. Intensions and Fregean compositionality
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## 3. Internal extensions

## 3. Internal extensions

## 3. Internal extensions

## GOAL

All
EXPRESSIONS have
external (or internal) semantic values $\langle\mathrm{X}\rangle, 《 \mathrm{X}\rangle, \ldots$ corresponding to (the contributions they make to) the functions of expressions in which they occur

## 3. Internal extensions

## GOAL

All
EXPRESSIONS have external or internal extensions 〈X〉 corresponding to the contributions they make to the external extensions of expressions in which they occur

## 3. Internal extensions

3. Internal extensions

Construction of internal extensions: standard method
3. Internal extensions

Construction of internal extensions:
standard method
3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.

## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X


## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
cf. Zimmermann (201I; 2012) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X


## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
cf. Zimmermann (2011; 2012) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
... i.e.: all values $\left\langle X_{i}+Y_{j}\right\rangle$ and $\left\langle Y_{j}\right\rangle$ have already been determined (externally, by previous applications of the standard method, or otherwise).


## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
cf. Zimmermann (2011; 2012) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$


## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
cf. Zimmerman (2011; 2012) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$
... ie.: whenever $\left\langle Y_{j}\right\rangle=\left\langle Y_{k}\right\rangle$, then:

$$
\left\langle X_{i}+Y_{j}\right\rangle=\left\langle X_{i}+Y_{k}\right\rangle .
$$

## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
cf. Zimmermann (2011; 2012) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$
-     + is representative for $X$


## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.

## Strategy

cf. Zimmermann (201I; 20I2) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$
-     + is representative for $X$
... i.e.: whenever $\left\langle X_{i}+Y_{k}\right\rangle=\left\langle X_{j}+Y_{k}\right\rangle$, for all $Y_{k}$, then:
$\left\langle Z\left[X_{i}\right]=Z\left[X_{i}\right]\right\rangle$, for all $Z[]$ already evaluated.


## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
cf. Zimmermann (2011; 2012) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$
-     + is representative for $X$


## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy cf. Zimmermann (201I; 20I2) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$
-     + is representative for $X$

Construction of $\left\langle X_{i}\right\rangle$

## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.

## Strategy

 cf. Zimmerman (201I; 20I2) for detailsTo extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$
-     + is representative for $X$

Construction of $\left\langle X_{i}\right\rangle$
$\left\langle X_{i}\right\rangle:=\lambda\left\langle Y_{j}\right\rangle .\left\langle X_{i}+Y_{j}\right\rangle$

## 3. Internal extensions

Construction of internal extensions: standard method

Cofinality assumption
Every expression occurs in some sentence.
Strategy
cf. Zimmerman (2011; 2012) for details
To extend the evaluation to a class $X$ of (valueless) expressions, choose a suitable construction +:

-     + is completed by X
-     + is compositional in $X$
-     + is representative for $X$

Construction of $\left\langle\mathrm{X}_{\mathrm{i}}\right\rangle$
$\left\langle\mathrm{X}_{\mathrm{i}}\right\rangle:=\lambda\left\langle\mathrm{Y}_{\mathrm{j}}\right\rangle .\left\langle\mathrm{X}_{i}+\mathrm{Y}_{\mathrm{j}}\right\rangle \quad \ldots$ and $\oplus$ is functional application.
3. Internal extensions

Construction of internal extensions:
standard method
3. Internal extensions

Construction of internal extensions: standard method

Example I (based on Ist approach to external extensions)
3. Internal extensions

Construction of internal extensions: standard method

Example I (based on Ist approach to external extensions)
$\langle$ is dead〉 $=\lambda\langle N P\rangle .\langle N P+$ is dead $\rangle$

## 3. Internal extensions

Construction of internal extensions: standard method

Example I (based on Ist approach to external extensions)
$\langle$ is dead $\rangle=\lambda\langle N P\rangle .\langle N P+$ is dead $\rangle$
(characteristic function of) set of individuals
3. Internal extensions

Construction of internal extensions: standard method

Example I (based on Ist approach to external extensions)
$\langle$ is dead〉 $=\lambda\langle N P\rangle .\langle N P+$ is dead $\rangle$

## 3. Internal extensions

Construction of internal extensions: standard method

Example I (based on Ist approach to external extensions)
$\langle$ is dead〉 $=\lambda\langle N P\rangle .\langle N P+$ is dead $\rangle$
$\cong$ external extension according to 2nd approach

## 3. Internal extensions

 Construction of internal extensions: standard methodExample I (based on Ist approach to external extensions)
$\langle$ is dead〉 $=\lambda\langle N P\rangle .\langle N P+$ is dead $\rangle$
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$\langle$ everybody $=\lambda$ VP＞．〈everybody＋VP〉
bound set variable！

## 3. Internal extensions

4 problems with standard method of constructing extensions
3. Internal extensions

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Problem I: Indeterminacy

3. Internal extensions

4 problems with standard method of constructing extensions

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Extensions (and other values) depend on choice of + .
3. Internal extensions

4 problems with standard method of constructing extensions

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Extensions (and other values) depend on choice of + .
However, the resulting value assignments (after completion) will always be isomorphic.
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However, the resulting value assignments (after completion) will always be isomorphic.

Solution:
Isomorphic theories should be declared notational variants of one another - provided they agree on the external extensions (and their 'interpretation').
3. Internal extensions

4 problems with standard method of constructing extensions

Problem 2: Laziness

3. Internal extensions 4 problems with standard method of constructing extensions

## Problem 2: Laziness

Internal extensions (and other values!) may still be in need of specification.
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3．Internal extensions
4 problems with standard method of constructing extensions

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E．g．，it is not obvious from
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that 〈 everybody〉 characterizes the supersets of〈person〉．

3．Internal extensions
4 problems with standard method of
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## Problem 2：Laziness

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E．g．，it is not obvious from
$\langle$ everybody〉 $=\lambda\langle V P\rangle$ ．〈everybody +VP$\rangle$ that 〈everybody〉 characterizes the supersets of〈person〉．
Solution：Background theory for characterizations of functional values．
3. Internal extensions

4 problems with standard method of constructing extensions

## Problem 3: Overdetermination

3. Internal extensions

4 problems with standard method of constructing extensions

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After internal extensions (...) have been constructed they may still appear in other constructions.
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Example
Heim \& Kratzer (1998)
$\langle V+$ everybody〉=???
3. Internal extensions

4 problems with standard method of constructing extensions

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Solution: No principled compositionality problems can arise (due to representativity of + ); however standards for specifying (functional) values are needed.
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## Example

Heim \& Kratzer (1998)

$$
\langle V+\text { everybody }\rangle=(\lambda x .\langle e v e r y b o d y\rangle(\lambda y .\langle V\rangle(y)(x))
$$

Solution: No principled compositionality problems can arise (due to representativity of + ); however standards for specifying (functional) values are needed.
3. Internal extensions

4 problems with standard method of constructing extensions

## Problem 4: No suitable + exists

3. Internal extensions

4 problems with standard method of constructing extensions

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Standard example
3. Internal extensions 4 problems with standard method of constructing extensions

## Problem 4: No suitable + exists

Standard example
3. Internal extensions

4 problems with standard method of constructing extensions

Problem 4: No suitable + exists
Standard example
Attitude verbs $\vee$; e.g. no $\oplus$ can satisfy:

$$
\langle V+S\rangle \neq\langle V\rangle \oplus\langle S\rangle
$$

because + is not compositional in V .
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## 4. Intensions and Fregean compositionality

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Solution strategy (for Problem 4): Local repair

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If + is not compositional in $X$, find alternative
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Solution strategy（for Problem 4）：Local repair
Frege（1892）
If + is not compositional in $X$ ，find alternative values 《Y》）and put：

$$
\left.\left\langle X_{i}\right\rangle:=\lambda 《 Y_{j}\right\rangle .\left\langle X_{i}+Y_{j}\right\rangle
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Attitude reports：
$\left\langle\right.$ believe：＝$\left.\lambda 《 S_{j}\right\rangle$ ．〈 believe $\left.+S_{j}\right\rangle$
．．．where（e．g．）《S》 is the intension of $S$

## 4. Intensions and Fregean compositionality

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## External intensions

## 4. Intensions and Fregean compositionality

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Carnap (1947), inspired by Wittgenstein (1922)

## 4. Intensions and Fregean compositionality

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Identify informational content with sets of possible worlds ('regions in Logical Space’), thereby obtaining intensions of (declarative sentences).

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$(!) 《 S\rangle \cong \lambda w .\langle S\rangle w$

## 4．Intensions and Fregean compositionality

## External intensions

Identify informational content with sets of possible worlds（＇regions in Logical Space’）， thereby obtaining intensions of（declarative sentences）．

Observation
$(!) 《 S\rangle \cong \lambda w$ ．〈S〉w
where $\langle S\rangle$ w is the extension of $S$ according to $w$ ．
Internal intensions
Generalize（！）from $S$ to arbitrary expressions．

## 4. Intensions and Fregean compositionality

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## Intensional compositionality

## 4．Intensions and Fregean compositionality

## Intensional compositionality

It is generally assumed that：

$$
《 X+Y 》=《 X\rangle \oplus \mathbb{Z}\rangle
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Fregean compositionality implies，but is not implied by， intensional compositionality．

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Laziness is particularly popular when it comes to applying Fregean compositionality．

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## Fregean Laziness

Laziness is particularly popular when it comes to applying Fregean compositionality. Zimmermann (1999), Larson (2002),... However

Relational analyses of referentially opaque transitive verbs are not the result of Fregean Laziness.
The compositional analysis
$(\lambda Y . \lambda X . Y \subseteq X)(M)\left(\lambda x . S\left(x,^{\wedge}(\lambda Y . \lambda X . Y \nless X)(W)\right)\right)$
of Every man seeks a woman.
can be obtained from the non-compositional modal paraphrase:
$(\forall x)\left[M(x) \rightarrow \square_{x}(\exists y)[W(y) \& L(x, y)]\right]$
by the standard method of constructing extensions.
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## 5. Outro

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## Thank you for your attention!

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